

# 符号から得られる多項式とデザイン理論の関係

田中 優帆 (Yuuho Tanaka)

Graduate School of Science and Engineering  
Waseda University

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- ▶ Preliminaries
- ▶ Designs and Molien series
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- ▶ Preliminaries
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# Combinatorial Designs

## Definition ( $t$ - $(n, k, \lambda)$ design)

Let

- ▶  $X := \{1, 2, \dots, n\}$
- ▶  $\mathcal{B} \subseteq \binom{X}{k}$

We say  $(X, \mathcal{B})$  is a  $t$ - $(n, k, \lambda)$  design, if there exists  $\lambda$  such that for all  $T \in \binom{X}{t}$ ,  $\lambda = |\{B \in \mathcal{B} \mid T \subseteq B\}|$ .

# Combinatorial Designs

## Example

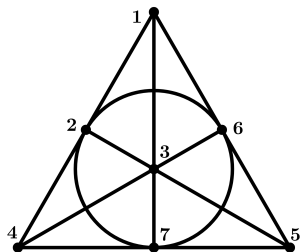
Let

$$\begin{cases} X = \{1, 2, 3, 4, 5, 6, 7\} \\ \mathcal{B} = \{\{1, 2, 4\}, \{2, 3, 5\}, \{3, 4, 6\}, \{4, 5, 7\}, \{5, 6, 1\}, \{6, 7, 2\}, \{7, 1, 3\}\} \end{cases}$$

Then,

$$\forall T \in \binom{X}{2}, |\{B \in \mathcal{B} \mid T \subseteq B\}| = 1.$$

$(X, \mathcal{B})$  is a 2-(7, 3, 1) design.



# Codes $\Rightarrow$ Combinatorial Designs

## Problem

*How do we obtain a  $t$ -designs structure?*

Let  $C$  be an  $\mathbb{F}_q$ -linear code of length  $n$ . For

$c = (c_1, c_2, \dots, c_n) \in C$ ,

- ▶  $\text{supp}(c) := \{i \mid c_i \neq 0\}$
- ▶  $\text{wt}(c) := |\text{supp}(c)|$
- ▶  $C_k := \{c \in C \mid \text{wt}(c) = k\}$

## Definition ( $t$ -design obtained from a code)

Let  $C$  be an  $\mathbb{F}_q$ -linear code of length  $n$ .

$$\begin{cases} X = \{1, 2, \dots, n\} \\ \mathcal{B}(C_k) = \{\text{supp}(c) \mid c \in C_k\} \subseteq \binom{X}{k} \end{cases}$$

If  $(X, \mathcal{B}(C_k))$  is a  $t$ - $(n, k, \lambda)$  design, then  $C_k$  is a  $t$ - $(n, k, \lambda)$  design.

We say  $C$  is a  $t$ -homogeneous if  $C_k$  is a  $t$ -design for all  $k$  ( $C_k \neq \emptyset$ ).

## Codes $\Rightarrow$ Combinatorial Designs

### Example (Hamming code)

Let  $H$  be the Hamming  $[7, 4]$  code (4-dim subspace of  $\mathbb{F}_2^7$ ) with the following generator matrix  $G$ :

$$G = \begin{bmatrix} 1 & 1 & 0 & 1 & 0 & 0 & 0 \\ 0 & 1 & 1 & 0 & 1 & 0 & 0 \\ 0 & 0 & 1 & 1 & 0 & 1 & 0 \\ 0 & 0 & 0 & 1 & 1 & 0 & 1 \end{bmatrix}$$

$$H_3 = \{(1, 1, 0, 1, 0, 0, 0), (0, 1, 1, 0, 1, 0, 0), (0, 0, 1, 1, 0, 1, 0), \\ (0, 0, 0, 1, 1, 0, 1), (1, 0, 0, 0, 1, 1, 0), (0, 1, 0, 0, 0, 1, 1), (1, 0, 1, 0, 0, 0, 1)\}, \\ \begin{cases} X = \{1, 2, 3, 4, 5, 6, 7\} \\ \mathcal{B}(H_3) = \{\{1, 2, 4\}, \{2, 3, 5\}, \{3, 4, 6\}, \{4, 5, 7\}, \{5, 6, 1\}, \{6, 7, 2\}, \{7, 1, 3\}\} \end{cases}$$

Then,

$$\forall T \in \binom{X}{2}, |\{B \in \mathcal{B} \mid T \subseteq B\}| = 1.$$

$(X, \mathcal{B}(H_3))$  is a  $2$ -( $7, 3, 1$ ) design.  $\Rightarrow H_3$  is a  $2$ -( $7, 3, 1$ ) design.

# Codes $\Rightarrow$ Combinatorial Designs

We can obtain a  $t$ -design from a code by using the following:

- ▶ Harmonic weight enumerators
- ▶ Jacobi polynomials (Main)



# Designs $\Leftrightarrow$ Harmonic weight enumerators

## Definition (Discrete harmonic function)

Let

- ▶  $\Omega = \{1, \dots, n\}$ ,  $X = 2^\Omega$ ,  $X_k = \binom{X}{k}$
- ▶  $\mathbb{R}X = \{\sum_{x \in X} c_i x \mid \forall i, c_i \in \mathbb{R}\}$
- ▶  $\mathbb{R}X_k = \{\sum_{x \in X_k} c_i x \mid \forall i, c_i \in \mathbb{R}\}$  ( $k = 0, 1, \dots, n$ )

For  $f \in \mathbb{R}X_k$ ,  $f$  is denoted by

$$f = \sum_{z \in X_k} f(z)z$$

and we can identify  $f$  with a function on  $X_k$  given by  $z \rightarrow f(z)$ .

Moreover,  $f$  can be extended to  $\tilde{f} \in \mathbb{R}X$ : for  $u \in X$ ,

$$\tilde{f}(u) = \sum_{z \in X_k, z \subset u} f(z).$$

## Designs $\Leftrightarrow$ Harmonic weight enumerators

- Let  $\gamma : \mathbb{R}X_k \rightarrow \mathbb{R}X_{k-1}$  be a map defined by linearity from

$$\gamma(z) = \sum_{y \in X_{k-1}, y \subset z} y \text{ for all } z \in X_k.$$

- $\text{Harm}_k = \ker(\gamma|_{\mathbb{R}X_k})$ .

Then, an element of  $\text{Harm}_k$  is a harmonic function.

### Example

Let  $\Omega = \{1, 2, 3, 4\}$  and

$$f = \{1, 2\} + \{1, 3\} - 2\{1, 4\} - 2\{2, 3\} + \{2, 4\} + \{3, 4\} \in \mathbb{R}X_2.$$

Then

$$\begin{aligned} \tilde{f}(\{1, 2, 3\}) &= f(\{1, 2\}) + f(\{1, 3\}) + f(\{2, 3\}) = 1 + 1 - 2 = 0, \\ \gamma(f) &= (\{1\} + \{2\}) + (\{1\} + \{3\}) \\ &\quad - 2(\{1\} + \{4\}) - 2(\{2\} + \{3\}) \\ &\quad + (\{2\} + \{4\}) + (\{3\} + \{4\}) = 0 \end{aligned}$$

and  $f \in \text{Harm}_2$ .

## Designs $\Leftrightarrow$ Harmonic weight enumerators

### Theorem (Delsarte ,1978)

Let  $C$  be a binary code of length  $n$ .

$C_\ell$  is a  $t$ -design  $\Leftrightarrow$

$$\sum_{c \in C_\ell} \tilde{f}(\text{supp}(c)) = 0, \forall f \in \text{Harm}_k \quad (1 \leq k \leq t)$$

### Example

Let

►  $C$  be a binary code of length 4.

►  $C_2 =$

$$\{(1, 1, 0, 0), (1, 0, 1, 0), (1, 0, 0, 1), (0, 1, 1, 0), (0, 1, 0, 1), (0, 0, 1, 1)\}.$$

►  $\mathcal{B}(C_2) = \{\{1, 2\}, \{1, 3\}, \{1, 4\}, \{2, 3\}, \{2, 4\}, \{3, 4\}\}.$

Then,  $C_2$  is a 2-(4, 2, 1) design.

$\Leftrightarrow$

$$f = \{1, 2\} + \{1, 3\} - 2\{1, 4\} - 2\{2, 3\} + \{2, 4\} + \{3, 4\} \in \text{Harm}_2.$$

$$\sum_{c \in C_2} \tilde{f}(\text{supp}(c)) = 1 + 1 - 2 - 2 + 1 + 1 = 0.$$

## Designs $\Leftrightarrow$ Harmonic weight enumerators

Definition (Harmonic weight enumerator (Bachoc, 1999))

A harmonic weight enumerator of  $C$  with  $f \in \text{Harm}_k$  is defined as follows:

$$w_{C,f}(x,y) = \sum_{c \in C} \tilde{f}(\text{supp}(c)) x^{n-\text{wt}(c)} y^{\text{wt}(c)}.$$

Example (Hamming code)

Let  $H$  be the Hamming  $[7, 4]$  code (4-dim subspace of  $\mathbb{F}_2^7$ ) with the following generator matrix  $G$ :

$$G = \begin{bmatrix} 1 & 1 & 0 & 1 & 0 & 0 & 0 \\ 0 & 1 & 1 & 0 & 1 & 0 & 0 \\ 0 & 0 & 1 & 1 & 0 & 1 & 0 \\ 0 & 0 & 0 & 1 & 1 & 0 & 1 \end{bmatrix}$$

$$H = \{(0, 0, 0, 0, 0, 0, 0), (1, 1, 0, 1, 0, 0, 0), (0, 1, 1, 0, 1, 0, 0), \dots\},$$

Then,

$$w_{H,f}(x,y) = \tilde{f}(\emptyset)x^{7-0}y^0 + \tilde{f}(\{1, 2, 4\})x^{7-3}y^3 + \tilde{f}(\{2, 3, 5\})x^{7-3}y^3 + \dots$$

# Designs $\Leftrightarrow$ Harmonic weight enumerators

Using

$$\begin{aligned}w_{C,f}(x,y) &= \sum_{c \in C} \tilde{f}(\text{supp}(c)) x^{n-\text{wt}(c)} y^{\text{wt}(c)} \\ &= \sum_{\ell=0}^n \left( \sum_{c \in C_\ell} \tilde{f}(\text{supp}(c)) \right) x^{n-\ell} y^\ell,\end{aligned}$$

we have the following theorem.

## Theorem

If for all  $f \in \text{Harm}_k$  ( $1 \leq k \leq t$ ),  $w_{C,f}(x,y) = 0$ , then, for all  $\ell$ ,  $C_\ell$  ( $\neq \emptyset$ ) is a  $t$ -design.

# Designs $\Leftrightarrow$ Harmonic weight enumerators

## Example (Hamming code)

$$\begin{aligned}w_{H,f}(x,y) &= \tilde{f}(\emptyset)x^7 + \tilde{f}(\{1,2,4\})x^4y^3 + \tilde{f}(\{2,3,5\})x^4y^3 + \dots \\ &= 0 + (f(\{1,2\}) + f(\{1,4\}) + f(\{2,4\}))x^4y^3 \\ &\quad + (f(\{2,3\}) + f(\{2,5\}) + f(\{3,5\}))x^4y^3 \\ &\quad + (f(\{1,3\}) + f(\{1,7\}) + f(\{3,7\}))x^4y^3 \\ &\quad + (f(\{1,5\}) + f(\{1,6\}) + f(\{5,6\}))x^4y^3 + \dots \\ &= 0 + (1 + 1 + 1)x^4y^3 + (1 + 1 + 1)x^4y^3 \\ &\quad + (1 - 5 - 5)x^4y^3 + (1 + 1 - 14)x^4y^3 + \dots \\ &= 0.\end{aligned}$$

$(f = \{1,2\} + \{1,3\} + \{1,4\} + \{1,5\} + \{1,6\} - 5\{1,7\} \cdots + 10\{6,7\} \in \text{Harm}_2)$

- ▶  $H_3$  is a 2-design.
- ▶  $H_4$  is a 2-design.

# Codes $\Rightarrow$ Combinatorial Designs

We can obtain a  $t$ -design from a code by using the following:

- ▶ Harmonic weight enumerators
- ▶ Jacobi polynomials (Main)

## Designs $\Leftrightarrow$ Jacobi polynomials

Let  $C$  be an  $\mathbb{F}_q$ -linear code of length  $n$ .

Definition (Weight enumerator)

$$W_C(x, y) := \sum_{c \in C} x^{n - \text{wt}(c)} y^{\text{wt}(c)}.$$

Definition (Jacobi polynomial (M. Ozeki, 1997))

Let  $T \subseteq \{1, \dots, n\} = [n]$ .

$$J_{C,T}(w, z, x, y) := \sum_{c \in C} w^{m_0(c)} z^{m_1(c)} x^{n_0(c)} y^{n_1(c)},$$

where

- ▶  $m_0(c) = |\{j \in T \mid c_j = 0\}|$
- ▶  $m_1(c) = |\{j \in T \mid c_j \neq 0\}|$
- ▶  $n_0(c) = |\{j \in [n] \setminus T \mid c_j = 0\}|$
- ▶  $n_1(c) = |\{j \in [n] \setminus T \mid c_j \neq 0\}|$



## Designs $\Leftrightarrow$ Jacobi polynomials

### Example (Hamming code)

Let  $H$  be the Hamming  $[7, 4]$  code (4-dim subspace of  $\mathbb{F}_2^7$ ) with the following generator matrix  $G$ :

$$G = \begin{bmatrix} 1 & 1 & 0 & 1 & 0 & 0 & 0 \\ 0 & 1 & 1 & 0 & 1 & 0 & 0 \\ 0 & 0 & 1 & 1 & 0 & 1 & 0 \\ 0 & 0 & 0 & 1 & 1 & 0 & 1 \end{bmatrix}$$

$$H = \{(0, 0, 0, 0, 0, 0, 0), (1, 1, 0, 1, 0, 0, 0), (0, 1, 1, 0, 1, 0, 0), \\ (0, 0, 1, 1, 0, 1, 0), (0, 0, 0, 1, 1, 0, 1), \dots, (1, 1, 1, 1, 1, 1, 1)\},$$

Then,

$$W_H(x, y) = x^{7-0}y^0 + x^{7-3}y^3 + \dots + x^{7-7}y^7 \\ = x^7 + 7x^4y^3 + 7x^3y^4 + y^7.$$

# Designs $\Leftrightarrow$ Jacobi polynomials

## Example (Hamming code)

$$H = \{(0, 0, 0, 0, 0, 0, 0), (1, 1, 0, 1, 0, 0, 0), (0, 1, 1, 0, 1, 0, 0), \\ (0, 0, 1, 1, 0, 1, 0), (0, 0, 0, 1, 1, 0, 1), \dots, (1, 1, 1, 1, 1, 1, 1)\},$$

$H_3$  is a  $2$ - $(7, 3, 1)$  design.

If  $T = \{1, 2\}$ , then

$$J_{H, \{1, 2\}}(w, z, x, y) \\ = w^2 z^0 x^5 y^0 + w^0 z^2 x^4 y^1 + \dots + w^0 z^2 x^0 y^5 \\ = w^2(x^5 + 2x^2y^3 + xy^4) + wz(4x^3y^2 + 4x^2y^3) + z^2(1x^4y + 2x^3y^2 + y^5).$$

The coefficient of  $z^t x^{n-k} y^{k-t}$  in  $J_{C, T} = |\{c \in C_k \mid T \subset \text{supp}(c)\}|$ .

## Theorem

If  $C$  is a  $t$ -homogeneous, then  $J_{C, T}$  is uniquely determined for all  $T \subset [n]$  ( $|T| = t$ ).

This Jacobi polynomial is denoted by  $J_{C, t}$ .

## Designs $\Leftrightarrow$ Jacobi polynomials

### Example (Hamming code)

Let  $H$  be the Hamming  $[7, 4]$  code (4-dim subspace of  $\mathbb{F}_2^7$ ) with the following generator matrix  $G$ :

$$G = \begin{bmatrix} 1 & 1 & 0 & 1 & 0 & 0 & 0 \\ 0 & 1 & 1 & 0 & 1 & 0 & 0 \\ 0 & 0 & 1 & 1 & 0 & 1 & 0 \\ 0 & 0 & 0 & 1 & 1 & 0 & 1 \end{bmatrix}$$

►  $H_3$  is a 2-(7, 3, 1) design.

►  $H_4$  is a 2-(7, 4, 2) design.

$\Rightarrow H$  is a 2-homogeneous.

$$J_{H,1}(w, z, x, y)$$

$$= w(x^6 + 4x^3y^3 + 3x^2y^4) + z(3x^4y^2 + 4x^3y^3 + y^6),$$

$$J_{H,2}(w, z, x, y)$$

$$= w^2(x^5 + 2x^2y^3 + xy^4) + wz(4x^3y^2 + 4x^2y^3) + z^2(x^4y + 2x^3y^2 + y^5).$$

# Contents

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- ▶ MacWilliams type identity

## Reference

- ▶ H. S. Chakraborty, T. Miezaki, M. Oura, Y. Tanaka. : Jacobi polynomials and designs theory I, *Discrete Mathematics*, **346** (2023) no. 6, No. 113339.

## Extended $t$ -design

Definition ( $t$ - $(n, k, \lambda_1^{a_1}, \dots, \lambda_N^{a_N})$  design)

Let

- ▶  $X := \{1, 2, \dots, n\}$
- ▶  $\mathcal{B} \subseteq \binom{X}{k}$

We say  $(X, \mathcal{B})$  is a  $t$ - $(n, k, \lambda_1^{a_1}, \dots, \lambda_N^{a_N})$  design, if there exists some  $\lambda_1, \dots, \lambda_N$ , where

$$a_i = |\{T \in \binom{X}{t} \mid \lambda_i = |\{B \in \mathcal{B} \mid T \subseteq B\}|\}| \quad (a_i > 0).$$

### Remark

If  $\lambda_1 \leq \dots \leq \lambda_N$ , then  $(X, \mathcal{B})$  is a

- ▶  $t$ - $(n, k, \lambda_N)$  packing design,
- ▶  $t$ - $(n, k, \lambda_1)$  covering design.

Moreover, the maximum (resp. minimum) number of blocks is denoted by  $D_{\lambda_N}(n, k, t)$  (resp.  $C_{\lambda_1}(n, k, t)$ ).

## Codes $\Rightarrow$ Combinatorial Designs

### Definition ( $t$ -design obtained from a code)

Let  $C$  be an  $\mathbb{F}_q$ -linear code of length  $n$ .

$$\begin{cases} X = \{1, 2, \dots, n\}, \\ \mathcal{B}(C_k) = \{\text{supp}(c) \mid c \in C_k\} \subseteq \binom{X}{k} \end{cases}$$

If  $(X, \mathcal{B}(C_k))$  is a  $t$ - $(n, k, \lambda_1^{a_1}, \dots, \lambda_N^{a_N})$  design, then  $C_k$  is a  $t$ - $(n, k, \lambda_1^{a_1}, \dots, \lambda_N^{a_N})$  design, where

$$a_i = |\{T \in \binom{X}{t} \mid \lambda_i = |\{B \in \mathcal{B}(C_k) \mid T \subseteq B\}|\}| \quad (\lambda_i > 0, a_i > 0).$$

When  $N = 1$ ,  $C_k$  is a  $t$ -design.

### Definition (Type III code)

A self-dual code  $C$  over  $\mathbb{F}_3$  of length  $n \equiv 0 \pmod{4}$  is called *Type III* if the weight of each codeword of  $C$  is a multiple of 3.

## Codes $\Rightarrow$ Combinatorial Designs

Example (Type III code of length 8  $C_{III}^8$  (unique))

$$X = \{1, 2, 3, 4, 5, 6, 7, 8\},$$

$$(C_{III}^8)_6 = \{(1, 2, 2, 0, 0, 2, 2, 2), (0, 1, 1, 1, 1, 2, 2, 0), (1, 1, 2, 1, 2, 0, 0, 1), (1, 2, 0, 1, 0, 2, 1, 2), (2, 0, 2, 2, 2, 2, 1, 0), \dots\},$$

$$\mathcal{B}((C_{III}^8)_6) = \{\{1, 2, 3, 6, 7, 8\}, \{2, 4, 5, 6, 7, 8\}, \{2, 3, 4, 5, 6, 7\}, \\ \{1, 2, 3, 4, 5, 7\}, \{1, 2, 3, 4, 5, 8\}, \{2, 3, 5, 6, 7, 8\}, \\ \{1, 2, 4, 6, 7, 8\}, \{2, 3, 4, 5, 6, 8\}, \{1, 2, 3, 5, 7, 8\}, \\ \{1, 3, 4, 5, 6, 8\}, \{1, 4, 5, 6, 7, 8\}, \{1, 3, 5, 6, 7, 8\}, \\ \{1, 2, 3, 4, 6, 8\}, \{1, 2, 3, 4, 6, 7\}, \{1, 2, 4, 5, 7, 8\}, \\ \{1, 3, 4, 5, 6, 7\}\},$$

Then,

$$a_1 = |\{T \in \binom{X}{2} \mid |\{B \in \mathcal{B}((C_{III}^8)_6) \mid T \subseteq B\}| = 8\}| \\ = |\{\{1, 2\}, \{1, 5\}, \{1, 6\}, \dots\}| = 12,$$

## Codes $\Rightarrow$ Combinatorial Designs

Example (Type III code of length 8  $C_{III}^8$  (unique))

$$\mathcal{B}((C_{III}^8)_6) = \{\{1, 2, 3, 6, 7, 8\}, \{2, 4, 5, 6, 7, 8\}, \{2, 3, 4, 5, 6, 7\}, \\ \{1, 2, 3, 4, 5, 7\}, \{1, 2, 3, 4, 5, 8\}, \{2, 3, 5, 6, 7, 8\}, \\ \{1, 2, 4, 6, 7, 8\}, \{2, 3, 4, 5, 6, 8\}, \{1, 2, 3, 5, 7, 8\}, \\ \{1, 3, 4, 5, 6, 8\}, \{1, 4, 5, 6, 7, 8\}, \{1, 3, 5, 6, 7, 8\}, \\ \{1, 2, 3, 4, 6, 8\}, \{1, 2, 3, 4, 6, 7\}, \{1, 2, 4, 5, 7, 8\}, \\ \{1, 3, 4, 5, 6, 7\}\},$$

Then,

$$a_2 = |\{T \in \binom{X}{2} \mid |\{B \in \mathcal{B}((C_{III}^8)_6) \mid T \subseteq B\}| = 9\}| \\ = |\{\{4, 5\}, \{1, 3\}, \{1, 4\}, \dots\}| = 16.$$

Then,  $(X, \mathcal{B}((C_{III}^8)_6))$  is a  $2-(8, 6, 8^{12}, 9^{16})$  design.

Therefore,  $(C_{III}^8)_6$  is a  $2-(8, 6, 8^{12}, 9^{16})$  design.

- ▶  $2-(8, 6, 9)$  packing design
- ▶  $2-(8, 6, 8)$  covering design



## Designs $\Leftrightarrow$ Jacobi polynomials

### Theorem

For any  $T \subseteq [n]$  ( $|T| = t$ ), there are precisely  $N$  members the coefficient of the term  $w^{i_1} z^{i_2} x^{j_1} y^{j_2}$  ( $i_2 + j_2 = k$ ) of Jacobi polynomials. Then,  $C_k$  is a  $t$ - $(n, k, \lambda_1^{a_1}, \dots, \lambda_N^{a_N})$  design.

### Example (Type II code of length 16)

Let  $C$  be the first Type II code of length 16 in “M. Harada and A. Munemasa, Database of self-dual codes”.

$C_{12}$  is a 2- $(16, 12, (15^{112}, 21^8))$  design.

$\Leftrightarrow$

$$J_{C,2}^1 = w^2(x^{14} + 15x^{10}y^4 + 47x^6y^8 + x^2y^{12}) + wz(12x^{11}y^3 + 104x^7y^7 + 12x^3y^{11}) + z^2(x^{12}y^2 + 47x^8y^6 + 15x^4y^{10} + y^{14}),$$

$$J_{C,2}^2 = w^2(x^{14} + 21x^{10}y^4 + 35x^6y^8 + 7x^2y^{12}) + 128wzx^7y^7 + z^2(7x^{12}y^2 + 35x^8y^6 + 21x^4y^{10} + y^{14}).$$

## Designs $\Leftrightarrow$ Jacobi polynomials

The number of the basis of  $J_{C,T} \Rightarrow$  We use "Molien series".

### Theorem

Jacobi polynomial  $J_{C_{III},T}$  is invariant under the action of group

$$G_3 = \left\langle \frac{1}{\sqrt{3}} \begin{bmatrix} 1 & 2 & 0 & 0 \\ 1 & -1 & 0 & 0 \\ 0 & 0 & 1 & 2 \\ 0 & 0 & 1 & -1 \end{bmatrix}, \begin{bmatrix} 1 & 0 & 0 & 0 \\ 0 & e^{2\pi i/3} & 0 & 0 \\ 0 & 0 & 1 & 0 \\ 0 & 0 & 0 & e^{2\pi i/3} \end{bmatrix} \right\rangle$$

of order 48.

# Designs $\Leftrightarrow$ Jacobi polynomials

Definition (Molien series)

$$M_{G_3}(u, v) = \sum_{n=0}^{\infty} \sum_{i=0}^n \dim M_{i, n-i} u^i v^{n-i}$$

is called *Molien series*, where

$$\begin{aligned} M_{i, n-i} &:= (\mathbb{C}[w, z, x, y]^{G_3})_{i, n-i}, \\ (\mathbb{C}[w, z, x, y]^{G_3})_{i, n-i} &= \{f \in \mathbb{C}[w, z, x, y] \mid \forall g \in G_3, gf = f, \\ &\quad \text{degree of } w, z \text{ in } f \text{ is } i, \\ &\quad \text{degree of } x, y \text{ in } f \text{ is } n - i\}. \end{aligned}$$

Theorem (R. P. Stanley, 1979)

$$M_{G_3}(u, v) = \frac{1}{|G_3|} \sum_{g \in G_3} \frac{1}{\det(1 - ug) \det(1 - vg)}.$$

## Designs $\Leftrightarrow$ Jacobi polynomials

### Example (Type III code of length 8 $C_{\text{III}}^8$ (unique))

The homogeneous part of degree 8 of the Molien series  $M_{G_3}(u, v)$  is

$$u^8 + u^7v + 2u^6v^2 + 2u^5v^3 + 2u^4v^4 + 2u^3v^5 + 2u^2v^6 + uv^7 + v^8.$$

Then,  $M_{2,6}$  is generated by 2 polynomials as follows:

$$M_{2,6} = \langle W_{C_{\text{III}}^4}(J_{C_{\text{III}}^4,2}), (J_{C_{\text{III}}^4,1})^2 \rangle$$

The coefficient of each basis is determined by comparing the terms of Jacobi polynomials.

$$J_{C_{\text{III}}^8,2}^1 = 1 \cdot W_{C_{\text{III}}^4}(J_{C_{\text{III}}^4,2}) + 0 \cdot (J_{C_{\text{III}}^4,1})^2$$

$$J_{C_{\text{III}}^8,2}^2 = 0 \cdot W_{C_{\text{III}}^4}(J_{C_{\text{III}}^4,2}) + 1 \cdot (J_{C_{\text{III}}^4,1})^2$$

We can obtain Jacobi polynomials as follows:

$$J_{C_{\text{III}}^8,2}^1 = w^2(x^6 + 8x^3y^3) + wz(4x^4y^2 + 32xy^5) + z^2(4x^5y + 32x^2y^4),$$

$$J_{C_{\text{III}}^8,2}^2 = w^2(4x^3y^3 + 4y^6) + wz(12x^4y^2 + 24xy^5) + 36z^2x^2y^4.$$

## Designs $\Leftrightarrow$ Jacobi polynomials

$$J_{C_{\text{III}},2}^1 = w^2(x^6 + 8x^3y^3) + wz(4x^4y^2 + 32xy^5) + z^2(4x^5y + 32x^2y^{6-2}),$$

$$J_{C_{\text{III}},2}^2 = w^2(4x^3y^3 + 4y^6) + wz(12x^4y^2 + 24xy^5) + z^2(36x^2y^{6-2}).$$

$(C_{\text{III}}^8)_6$  is a  $2$ -( $8, 6, (8^{12}, 9^{16})$ ) design.

$$\begin{aligned}\Rightarrow \lambda_1 = 8 &= \frac{1}{4} \times (\text{the coefficient of the term } z^2 x^2 y^{6-2} \text{ of } J_{C_{\text{III}},2}^1) \\ &= \frac{1}{4} \times 32,\end{aligned}$$

$$\begin{aligned}\lambda_2 = 9 &= \frac{1}{4} \times (\text{the coefficient of the term } z^2 x^2 y^{6-2} \text{ of } J_{C_{\text{III}},2}^2) \\ &= \frac{1}{4} \times 36.\end{aligned}$$

III	$2$ -( $8, 6, (8^{12}, 9^{16})$ )
	$4$ -( $16, 9, (88^{28}, 92^{336}, 94^{896}, 96^{560})$ )
	$2$ -( $20, 9, (392^{90}, 432^{100})$ )
	$2$ -( $20, 12, (4212^{100}, 4296^{90})$ )
	$2$ -( $20, 15, (6308^{90}, 6444^{100})$ )
IV	$2$ -( $6, 4, (1^{12}, 2^3)$ )

## Ex: Type III code of length 8 $C_{\text{III}}^8$ (unique)

$$W_{C_{\text{III}}^8}(x, y) = x^8 + 16x^5y^3 + 64x^2y^6.$$

$(C_{\text{III}}^8)_6$  is a 2-(8, 6, 8) packing design and a 2-(8, 6, 9) covering design.

$$C_8(8, 6, 2) \leq 16 \leq D_9(8, 6, 2)$$

$$\begin{aligned} \Rightarrow 16 &= 1/4 \times (\text{the coefficient of the term } x^2y^6 \text{ of } W_{C_{\text{III}}^8}) \\ &= 1/4 \times 64, \end{aligned}$$

III	$C_8(8, 6, 2) \leq 16 \leq D_9(8, 6, 2)$
	$C_{88}(16, 9, 4) \leq 1360 \leq D_{96}(16, 9, 4)$
	$C_{392}(20, 9, 2) \leq 2180 \leq D_{432}(20, 9, 2)$
	$C_{4212}(20, 12, 2) \leq 12240 \leq D_{4296}(20, 12, 2)$
	$C_{6308}(20, 15, 2) \leq 11544 \leq D_{6444}(20, 15, 2)$
IV	$C_1(6, 4, 2) \leq 3 \leq D_2(6, 4, 2)$

# Contents

- ▶ Preliminaries
- ▶ Designs and Molien series
- ▶ MacWilliams type identity

## Reference

- ▶ H. S. Chakraborty, T. Miezaki, M. Oura, Y. Tanaka. : Jacobi polynomials and designs theory I, *Discrete Mathematics*, **346** (2023) no. 6, No. 113339.

## MacWilliams type identity

### Theorem (M. Ozeki, 1997)

If  $C$  is a binary linear code, then

$$J_{C^\perp, T}(w, z, x, y) = \frac{1}{|C|} \left( \begin{bmatrix} 1 & 1 \\ 1 & -1 \end{bmatrix} \otimes I \right) J_{C, T}(w, z, x, y).$$

### Example (Hamming code)

Let  $H^\perp$  be the dual code of  $H$ .

Then,

$$\begin{aligned} J_{H^\perp, 1}(w, z, x, y) &= w(x^6 + 3x^2y^4) + 4zx^3y^3 \\ &= J_{H, 1}(w + z, w - z, x + y, x - y), \\ J_{H^\perp, 2}(w, z, x, y) &= w^2(x^5 + xy^4) + 4wzx^2y^3 + 2z^2x^3y^2 \\ &= J_{H, 2}(w + z, w - z, x + y, x - y). \end{aligned}$$

$$\ast J_{H, 1}(w, z, x, y) = w(x^6 + 4x^3y^3 + 3x^2y^4) + z(3x^4y^2 + 4x^3y^3 + y^6),$$

$$J_{H, 2}(w, z, x, y) = w^2(x^5 + 2x^2y^3 + xy^4) + wz(4x^3y^2 + 4x^2y^3) + z^2(x^4y + 2x^3y^2 + y^5).$$



# Jacobi polynomial with respect to $\ell$ reference vectors

## Definition (Jacobi polynomial)

Let  $C$  be an  $\mathbb{F}_q$ -linear code of length  $n$ . Let  $\mathbf{T} := (T_1, \dots, T_\ell)$  of pairwise disjoint sets  $T_i \subseteq [n]$ .

$$J_{C, \mathbf{T}}(w_1, z_1, \dots, w_\ell, z_\ell, x_0, x_1) := \sum_{c \in C} w_1^{m_0^1(c)} z_1^{m_1^1(c)} \dots w_\ell^{m_0^\ell(c)} z_\ell^{m_1^\ell(c)} x_0^{n_0(c)} x_1^{n_1(c)},$$

where

- ▶  $m_0^i(c) = |\{j \in T_i \mid c_j = 0\}|$  ( $i = 1, \dots, \ell$ )
- ▶  $m_1^i(c) = |\{j \in T_i \mid c_j \neq 0\}|$  ( $i = 1, \dots, \ell$ )
- ▶  $n_0(c) = |\{j \in [n] \setminus \bigcup_{1 \leq i \leq \ell} T_i \mid c_j = 0\}|$
- ▶  $n_1(c) = |\{j \in [n] \setminus \bigcup_{1 \leq i \leq \ell} T_i \mid c_j \neq 0\}|$

## Jacobi polynomial with respect to $\ell$ reference vectors

Example (Type III code of length 4  $C_{\text{III}}^4$  (unique))

$$C_{\text{III}}^4 = \{(0, 0, 0, 0), (1, 0, 1, 1), (2, 0, 2, 2), (2, 1, 0, 1), \\ (0, 1, 1, 2), (1, 1, 2, 0), (1, 2, 0, 2), (2, 2, 1, 0), (0, 2, 2, 1)\}$$

If  $\mathbf{T} = (\{1, 2\}, \{4\})$  then

$$\begin{aligned} J_{C_{\text{III}}^4, (\{1, 2\}, \{4\})}(w_1, z_1, w_2, z_2, x_0, x_1) \\ &= w_1^2 z_1^0 w_2^1 z_2^0 x_0^1 x_1^0 + \cdots + w_1^1 z_1^1 w_2^0 z_2^1 x_0^0 x_1^1 \\ &= w_1^2 w_2 x_0 + 4w_1 z_1 z_2 x_1 + 2z_1^2 w_2 x_1 + 2z_1^2 z_2 x_0. \end{aligned}$$

## Jacobi polynomial with respect to $\ell$ reference vectors

Example (Type III code of length 4  $C_{\text{III}}^4$  (unique))

If  $\mathbf{T} = (\{1, 2\}, \{4\})$  then

$$\begin{aligned} & J_{C_{\text{III}}^4, (\{1, 2\}, \{4\})}(w_1, z_1, w_2, z_2, x_0, x_1) \\ &= w_1^2 w_2 x_0 + 4w_1 z_1 z_2 x_1 + 2z_1^2 w_2 x_1 + 2z_1^2 z_2 x_0. \end{aligned}$$

Then,

$$\begin{aligned} & J_{(C_{\text{III}}^4)^\perp, (\{1, 2\}, \{4\})}(w_1, z_1, w_2, z_2, x_0, x_1) \\ &= \frac{1}{9} J_{C_{\text{III}}^4, (\{1, 2\}, \{4\})}(w_1 + 2z_1, w_1 - z_1, w_2 + 2z_2, w_2 - z_2, x_0 + 2x_1, x_0 - x_1) \\ &= \frac{1}{|C_{\text{III}}^4|} \left( \begin{bmatrix} 1 & 2 \\ 1 & -1 \end{bmatrix} \otimes I \otimes I \right) J_{C_{\text{III}}^4, (\{1, 2\}, \{4\})}(w_1, z_1, w_2, z_2, x_0, x_1). \end{aligned}$$

We can obtain a MacWilliams type identity for the Jacobi polynomial of  $C_{\text{III}}^4$  with respect to 2 reference vectors.

## MacWilliams type identity

We can obtain a MacWilliams type identity for the Jacobi polynomial of an  $\mathbb{F}_q$ -linear code  $C$  with respect to  $\ell$  reference vectors as follows:

Theorem (Y. Tanaka et al., 2023)










$$J_{C^\perp, \mathbf{T}}(w_1, z_1, \dots, w_\ell, z_\ell, x_0, x_1) = \frac{1}{|C|} \left( \begin{bmatrix} 1 & q-1 \\ 1 & -1 \end{bmatrix} \otimes I \otimes \dots \otimes I \right) J_{C, \mathbf{T}}(w_1, z_1, \dots, w_\ell, z_\ell, x_0, x_1).$$

Theorem (M. Ozeki, 1997)

If  $C$  is a binary linear code, then

$$J_{C^\perp, T}(w, z, x, y) = \frac{1}{|C|} \left( \begin{bmatrix} 1 & 1 \\ 1 & -1 \end{bmatrix} \otimes I \right) J_{C, T}(w, z, x, y).$$

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